Chapter MK:V

V. Diagnoseansätze

- Diagnoseproblemstellung
- Diagnose mit Bayes
- □ Evidenztheorie von Dempster/Shafer
- Diagnose mit Dempster/Shafer
- Truth Maintenance
- Assumption-Based TMS
- Diagnosis Setting
- Diagnosis with the GDE
- Diagnosis with Reiter
- □ Grundlagen fallbasierten Schließens
- □ Fallbasierte Diagnose

Non-Monotonicity

Situations where reasoning becomes non-monotonic:

- □ dynamic worlds, where "facts" can change
- observed input depends on time and place
- □ new input causes retractions, e.g., a hypothesis is discovered to be false
- exceptions become known

Foundations:

1. Deduction

Collection: Logics, Part: Propositional Logics

2. Non-Monotonicity

Collection: Logics, Part: Extensions

- □ The examples are typical for many diagnosis situations.
- □ If any of the mentioned situations occurs, one can simply restart the reasoning process from scratch in order to determine the valid deductions.
- A reason (or truth) maintenance system is some kind of bookkeeping mechanism that keeps track whether a fact is still inferable or not. The rationale is that the organizational overhead for bookkeeping is often smaller than an entire restart of the reasoning process.
- □ Truth maintenance system (TMS) in fact is a misnomer; rather say reason maintenance system (RMS).

Operationalization

Idea: Identification and retraction of inferences that are no longer valid.

Well-known TMS concepts:

- dependency-directed backtracking
- justification-based truth maintenance (JTMS)
- □ assumption-based truth maintenance (ATMS)

Interplay between inference engine and TMS:



Definition 12 (Datum, Node)

A datum is the smallest unit of information we are interested in. A node is the data structure of a TMS to represent a datum. There is a one-to-one correspondence between nodes and datums.

Illustration:



Examples for a datum:



Communication between the TMS and the inference engine is in terms of nodes:



Nodes are interpreted differently:

□ Inference engine.

Perform deductions that are grounded on the datum's semantics.

□ TMS.

Perform deductions respecting a node's validity. These deductions are grounded on logical implications.

Example for an inference engine deduction:

- N0001: IF student(X) THEN underpaid(X) AND overworked(x)
- N0002: student(robbie)
- N0003: underpaid(robbie) AND overworked(robbie)

- A datum is worth to be to be remembered / maintained / reasoned about / retracted / assumed. A datum can be an assertion, a fact, an inference rule, a procedures, a computation result, etc.
- □ The TMS cannot perform deductions that are grounded on the datum's semantics—but, the inference engine can inform the TMS on a deduction like N0001 ∧ N0002 → N0003.

The inference engine communicates the important (a subset of all) deductions to the TMS by means of justifications.

Definition 13 (Justification)

A justification is a condition on a node and consists of three parts:

- 1. Consequent. The node of the datum that has been inferred by the inference engine.
- 2. Antecedents list. Nodes of the datums that are used by the inference engine to infer the datum of the consequent.
- 3. Informant. A comment that explains the inference in more detail.

Example for an inference engine deduction:

- N0001: IF student(X) THEN underpaid(X) AND overworked(x)
- N0002: student(robbie)
- N0003: underpaid(robbie) AND overworked(robbie)

Justification communicated to the TMS:

 \langle N0003, MODUS-PONENS, {N0001, N0002} \rangle

- □ In the example, "N0003" is the consequent, "MODUS-PONENS" is the informant, and "{N0001, N0002}" are the antecedents.
- □ The example can be considered as a formulation in predicate logics.

Definition 14 (TMS Node Types)

TMS nodes can be of the following types:

1. Premise.

The *inference engine* has indicated that the associated datum is always true.

2. Contradiction.

The *inference engine* has indicated that the associated datum is always false.

3. Assumption.

The *inference engine* has indicated that the associated datum is true unless stated otherwise.

4. Justified.

Based on premise nodes, assumption nodes, and justifications, the node can be logically inferred by the *TMS*.

- □ With each node a label is associated; it encodes the current truth value (belief) of its node.
- □ Assumption nodes are used to represent datums which the inference engine has chosen to believe at the moment, but which it may want to retract later.
- □ The associated datum of a justified node is true.
- □ The TMS informs the inference engine if a node is justified or not, and if a contradiction can be inferred.
- Depending on the associated node's type, the colloquial semantics of the datum "student(robbie)" is as follows:
 - Robbie is a student for ever (premise).
 - Robbie is not student (contradiction).
 - Let's say, Robbie is a student (assumption).
 - The TMS can infer that Robbie is a student (justified).

Graphical notation, adopted from [Forbus/deKleer 1993]:



How Justifications Help

Identifying responsibilities for conclusions: Why does Z hold?



How Justifications Help (continued)

Guiding backtracking: Which assumption shall be retracted?



Formalization

Specification of TMS concepts in propositional logics:

- □ Every TMS node can be interpreted as a propositional symbol.
- □ Every justification can be interpreted as a propositional definite Horn clause. If the nodes x_1, \ldots, x_m justify node *n*, this is represented as:

$$\neg x_1 \lor \ldots \lor \neg x_m \lor n \quad \approx \quad (x_1 \land \ldots \land x_m) \to n$$

- \Box A premise node *n* corresponds to a unit clause *n*.
- \Box A contradiction node can be specified by the negative unit clause $\neg n$.

The truth value (belief) of a node is stored in its label; it can be encoded in different ways.

Concepts of a JTMS:

- A node can be labeled either as : IN or as : OUT and indicates whether or not it can be logically inferred from premise nodes, assumption nodes, and justifications.
- The JTMS answers queries respecting a node's truth value by simply returning its label.

Problems when using a JTMS:

- □ Assumption changes require relabeling.
- Cyclic justifications must be detected and treated specifically.

- □ If the rate of assumption changes is much larger then the number of queries about node labels, much time is wasted in relabeling.
- □ The JTMS is unsuitable as truth maintenance concept if the problem solving task requires frequent context switches or the parallel analysis of several contexts.

Label (continued)

Concepts of an ATMS:

- The label of a node n contains sets of assumption nodes. Using the given premises and justifications, n can be logically inferred from each set of assumption nodes.
- □ The ATMS answers queries respecting a node's truth value by checking the node's label. Example: Can node *n* be inferred from the assumptions $\{A, B, C\}$?
- If a contradiction node can be inferred from a set of assumptions, this set is removed from all labels.

Problem when using an ATMS:

The size of a node's label, i. e., the number of sets of assumptions where this node holds in, can grow exponentially in the number of assumption nodes.

- □ The inference engine should introduce as few assumptions as possible.
- The ATMS ensures that no node follows from a set of assumptions if a contradiction node also follows from that set. This implies that no special contradiction handling is necessary for the ATMS.

JTMS.

A context switch means:

- 1. Relabel assumptions respective desired : IN or : OUT value.
- 2. Retract justifications that are no longer valid.
- 3. Propagate justifications that are valid now.
- → JTMS incurs the cost of context switching when the actual context switches are made.

ATMS.

Context switches are free since all reasonable contexts are constructed already.

- A "reasonable context" is a context from which something can be inferred.
 - → ATMS incurs the cost of context switching up front.

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Assumption-Based TMS (ATMS) Example



- **\Box** Assumptions made: {*A*, *B*, *D*, *E*}
- **□** Result: *h* follows from $\{A, B, D, E\}$

Example (continued):



- **D** Retract: A, E
- **\Box** Result: h follows from $\{B, D\}$

In particular, h follows from: $\{A, D\}, \{A, B, D\}, \{A, D, E\}, \{A, B, D, E\}, \{B, D\}, \{B, D, E\}$

Terminology

D Environment.

A set of assumption nodes, short: assumptions.

□ Holds / is supported.

A node holds in / is supported by an environment E, if there exists a set of justifications such that the node can be logically inferred from E.

□ Label (of a node).

A label of a node comprises environments where the node holds in.

□ Nogood set.

An environment in which some contradiction node holds.

Consistent.

A consistent environment is one which is not a nogood set.

□ Context (of an environment).

A context of an environment is the set of nodes that hold in the environment.

- \Box The node *h* holds in six environments.
- **D** The context of the assumptions $\{A, B\}$ is $\{r, g\}$.

ATMS Labels

Objective: Determine whether a node n holds in some environment.

Naive solution: Record all environments in which n holds.

Better: Exploit monotonicity: If n follows from some environment E, it will follow from any superset of E as well. Discard nogood sets from a node's label.



ATMS Labels (continued)



- $\Box \quad \text{The label of } z \text{ is } \langle z, \{\{S, T\}\} \rangle.$
- \square {*R*, *S*} seems to support node *z*, but is not included in the label because it is a nogood set. More precisely:
 - $\{R, S\}$ contains an environment, $\{R\}$, which is a nogood set.

ATMS Labels (continued)

Following labels play a special role:

(a) Label with no environment: $\langle n, \{\} \rangle$

- □ The dependency network of justifications contains no pathway from the set of assumptions and premises to the node *n*, or
- □ all potential label environments are nogood sets.

(b) Label with the empty environment: $\langle n, \{\{\}\} \rangle$

 \Box The node *n* holds in every environment.

- \Box If a node *n* has a label with no environment then *n* has no support; it does not mean that *n* is a contradiction.
- □ All premises have a label consisting of the empty environment.
- Non-premise nodes can hold in empty environments if they ultimately depend only on premises.

Summary of ATMS concepts in propositional logics:

- **\Box** The ATMS nodes define a set of propositional symbols Σ .
- \Box A subset $A \subset \Sigma$ of symbols are marked as assumptions.
- □ An ATMS justification is encoded as a definite Horn clause.
- \Box If *n* is a premise node, it is represented as unit clause *n*.
- □ If *n* is a contradiction node, it is represented as the negative unit clause $\neg n$.
- $\hfill\square$ An environment ${\bf E}$ is a subset of assumptions, ${\bf E} \subset {\bf A}.$

Definition 15 (Hold, Nogood Set, Miminality)

Let α be the logical conjunction of premises and justifications, and let α be consistent. Moreover, let β be the logical conjunction of the assumptions of an environment E.

- 1. A node *n* is said to hold in the environment **E**, if $\alpha \land \beta \models n$.
- 2. If $\alpha \land \beta \models n$ and if *n* is a contradiction node, then the environment **E** is a nogood set.
- 3. A nogood set is minimum if it contains no other nogood set as a subset.

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ATMS recording service:

The ATMS receives a stream of nodes, assumptions, and justifications.

ATMS information service:

The ATMS answers queries concerning environments and nodes that may hold in these environments.

Definition 16 (ATMS Label Properties)

To facilitate answering queries, the ATMS maintains for a label $\{E_1, \ldots, E_k\}$ of some node *n* the following four properties:

1. Soundness.

n holds in each \mathbf{E}_i .

2. Consistency.

No contradiction node can be derived from some $\alpha \wedge \beta_i$. α is the conjunction of recorded premises and justifications, and β_i is be the conjunction of the assumptions in \mathbf{E}_i .

3. Completeness.

Every consistent environment *E* in which *n* holds is a superset of some E_i .

4. Minimality.

No E_i is a proper subset of any other.

Label Update

Concepts:

- □ Label updating happens incremental.
- Labels are made locally correct and label changes are propagated until labels become globally correct.
- Assumptions are created with labels containing the single environment containing themselves.
- □ All other nodes are created with empty labels.

Label Update (continued)



Let L_{i_j} be the label of the *i*th node of the *j*th justification for some node *n*. Label update algorithm for node *n*:

- 1. Compute a tentative label $L' = \{ \bigcup e_i \mid e_i \in L_{i_j} \}$
- 2. Remove from L' all nogood sets, and all environments subsumed by other environments in L'.

Label Update (continued)

Example:

- $\Box \quad \langle e, \{\{A,B\}, \{C\}\} \rangle$
- $\square \ \langle f, \{\{A\}, \{D\}\} \rangle$
- $\square \ \langle \bot, \{\{C, D\}\}\rangle$

New justification j from inference engine: $e \wedge f \rightarrow g$

Label Update (continued)

Example:

 $\Box \quad \langle e, \{\{A, B\}, \{C\}\} \rangle$ $\Box \quad \langle f, \{\{A\}, \{D\}\} \rangle$ $\Box \quad \langle f, \{\{A\}, \{D\}\} \rangle$

 $\Box \quad \langle \bot, \{\{C, D\}\}\rangle$

New justification j from inference engine: $e \wedge f \rightarrow g$

Computation of a new label for node g:

$$\begin{array}{l} & \\ \mathsf{Subsumes} \\ \mathsf{Subsumes} \\ \mathsf{Subsumes} \\ \end{aligned}$$

Label Update (continued)

Global label update algorithm:

- 1. To update node n, compute its new label L' as described before.
- 2. If the label is not changed, return.
- 3. If n is a contradiction node:
 - (a) Mark all environments of L' as nogood sets.
 - (b) Remove all new nogood sets from every node label.
- 4. If *n* is not a contradiction node, then recursively update all the consequences of *n*.

□ This algorithm is inefficient (cf. Step 1). A more efficient version propagates merely incremental changes of node labels.